

On the computation of A_∞ -maps

A. Berciano, **M.J. Jiménez**, P. Real

Universidad de Sevilla

September 2007

Summary

Notion of A_∞ -structures.

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Theoretical results to compute them.

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Theoretical results to compute them.

Improvements in the theoretical algorithms.

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- 1 Introduction
 - Preliminaries
 - Historical origin
 - Mathematical notion
- 2 A_∞ -structures via perturbation
 - Contractions
 - The tensor trick
- 3 Theoretical study of complexity
 - Main Results
 - Computational advantages

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Rings

Let us suppose that Λ is a commutative ring with $1 \neq 0$.

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Modules

A *DG-module* is a graded module $M = \{M_n\}_{n \geq 0}$, endowed with a differential $d : M \rightarrow M$ (that is, a morphism of graded modules of degree -1 such that $d^2 = 0$).

Algebras

Definition

A **differential graded algebra** (A, μ_A, η) , or simply DG-algebra, is a DG-module equipped with two morphisms $\mu_A : A \otimes A \rightarrow A$ and $\eta : \Lambda \rightarrow A$, such that μ_A is an **associative product**, i.e., $\mu_A(\mu_A \otimes 1) = \mu_A(1 \otimes \mu_A)$, and η is a bilateral unit $\eta : \Lambda \rightarrow A$, i.e.,

$$\begin{array}{ccccc}
 \Lambda \otimes A & \xrightarrow{\cong} & A & \xrightarrow{\cong} & A \otimes \Lambda \\
 \downarrow \eta \otimes 1 & & \parallel & & \downarrow 1 \otimes \eta \\
 A \otimes A & \xrightarrow{\mu_A} & A & \xleftarrow{\mu_A} & A \otimes A
 \end{array}$$

A_∞ -structures: Historical evolution

The notion of A_∞ -algebras appears in the literature as a generalization of “**associative up to homotopy**”.

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- there exists $\mu_3 : M^{\otimes 3} \rightarrow M$ of degree $+1$, such that

$$\mu_2(\mu_2 \otimes 1) - \mu_2(1 \otimes \mu_2).$$

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$$\mu_3 d + d \mu_3 = \mu_2(\mu_2 \otimes 1) - \mu_2(1 \otimes \mu_2) \neq 0.$$

A_∞ -structures: Classical Example

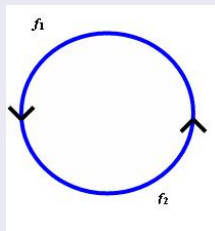
Let $(X, *)$ be a topological space with a base point $*$ and let ΩX denote the **space of based loops in X** : a point of ΩX is a continuous map $f : \mathbb{S}^1 \rightarrow X$ taking the base point of the circle to the base point $*$.

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Let us take as **multiplication** the composition map

$$\begin{aligned} \mu_2 : \Omega X \times \Omega X &\rightarrow \Omega X \\ (f_1, f_2) &\longrightarrow f_1 * f_2 \end{aligned}$$

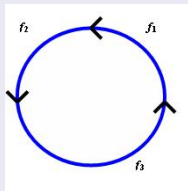


A_∞ -structures: Classical Example

Non associative

μ_2 is **not associative** because of

$$(f_1 * f_2) * f_3$$

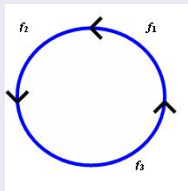


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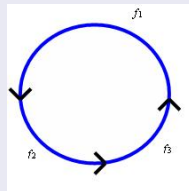
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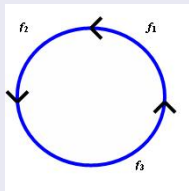


A_∞ -structures: Classical Example

Non associative

 μ_2 is associative up to homotopy

$$(f_1 * f_2) * f_3$$

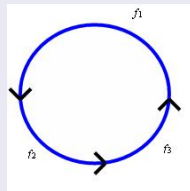


Homotopy operator



$$\mu_3 : [0, 1] \times (\Omega X)^3 \rightarrow \Omega X$$

$$f_1 * (f_2 * f_3)$$



A_∞ -structures: Mathematical notion

Definition

An A_∞ -algebra is a graded module A , with a family of graded maps $m_i : A^{\otimes i} \rightarrow A$, of degree $i - 2$, such that for all $i \geq 1$:

$$\sum_{n=1}^i \sum_{k=0}^{i-n} (-1)^{n+k+nk} m_{i-n+1}(1^{\otimes k} \otimes m_n \otimes 1^{\otimes i-n-k}) = 0.$$

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- $m_3(m_1 \otimes 1^2 + 1^2 \otimes m_1 + 1 \otimes m_1 \otimes 1) + m_1 m_3 = m_2(m_2 \otimes 1 - 1 \otimes m_2)$.

A_∞ -structures: Mathematical notion

Example

Every **DG-algebra** is, in particular, an A_∞ -algebra with $m_1 = d$, $m_2 = \mu$ and $m_i = 0$ for all $i \geq 3$.

A_∞ -structures: Mathematical notion

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Every **DG-algebra** is, in particular, an A_∞ -algebra with $m_1 = d$, $m_2 = \mu$ and $m_i = 0$ for all $i \geq 3$.

Example

The chain complex of the loop space of X , $C_*(\Omega X)$ has an A_∞ -algebra structure.

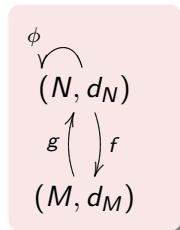
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Contractions

Definition

A *contraction* $c : \{N, M, f, g, \phi\}$ is a 5-tuple, such that

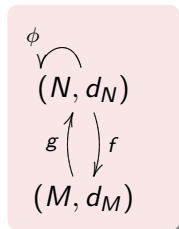


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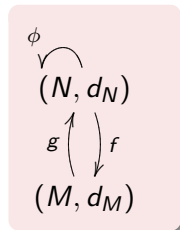


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- $f : N_* \rightarrow M_*$ and $g : M_* \rightarrow N_*$ morphisms of degree zero;

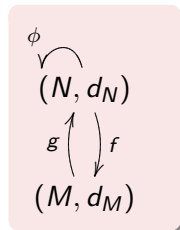


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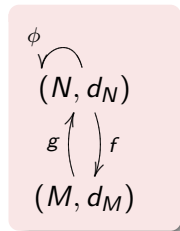


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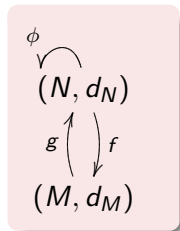


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- $f\phi = 0, \quad \phi g = 0, \quad \phi\phi = 0$.



Contractions: A_∞ -structures

Theorem (Kad80, GLS91)

Let (A, d_A, μ) and (M, d_M) be a connected DG-algebra and a DG-module, respectively and $c : \{A, M, f, g, \phi\}$ a contraction between them. Then the DG-module M is provided with an A_∞ -algebra structure

Contractions: A_∞ -structures

given by

$$m_n : M^{\otimes n} \rightarrow M$$

$$m_1 = -d_M$$

$$m_n = (-1)^{n+1} f \mu^{(1)} \phi^{[\otimes 2]} \mu^{(2)} \dots \phi^{[\otimes n-1]} \mu^{(n-1)} g^{\otimes n}, \quad n \geq 2 \quad (1)$$

with

$$\mu^{(k)} = \sum_{i=0}^{k-1} (-1)^{i+1} 1^{\otimes i} \otimes \mu_A \otimes 1^{\otimes k-i-1},$$

$$\phi^{[\otimes k]} = \sum_{i=0}^{k-1} 1^{\otimes i} \otimes \phi \otimes (g f)^{\otimes k-i-1}$$

Contractions: A_∞ -structures

Computational Consequence

A contraction from a DG-algebra A to a DG-module M provides an algorithm to compute an A_∞ -algebra structure on M .

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Main Results

Theorem

Any composition of the kind $\phi^{[\otimes s]} \mu^{(s)}$ ($s = 2, \dots, n-1$) in the formula (1), which is given by

$$\left(\sum_{j=0}^{s-1} 1^{\otimes j} \otimes \phi \otimes (g f)^{\otimes s-j-1} \right) \circ \left(\sum_{i=0}^{s-1} (-1)^{i+1} 1^{\otimes i} \otimes \mu_A \otimes 1^{\otimes s-i-1} \right),$$

can be reduced to

$$\phi^{[\otimes s]} \mu^{(s)} = \sum_{i=0}^{s-1} (-1)^{i+1} 1^{\otimes i} \otimes \phi \mu_A \otimes 1^{\otimes s-i-1}.$$

Main Results

$$\phi^{[\otimes s]} \mu^{(s)} = \sum_{i=0}^{s-1} (-1)^{i+1} 1^{\otimes i} \otimes \phi \mu_A \otimes 1^{\otimes s-i-1}. \quad (2)$$

Main Results

$$\phi^{[\otimes s]} \mu^{(s)} = \sum_{i=0}^{s-1} (-1)^{i+1} 1^{\otimes i} \otimes \phi \mu_A \otimes 1^{\otimes s-i-1}. \quad (2)$$

Moreover, given a composition of the kind

$$(\phi^{[\otimes s-1]} \mu^{(s-1)}) \circ (\phi^{[\otimes s]} \mu^{(s)}) \quad s = 3, \dots, n-2,$$

for every index i in the sum (2) of $\phi^{[\otimes s]} \mu^{(s)}$, the formula of $\phi^{[\otimes s-1]} \mu^{(s-1)}$ in such a composition can be reduced to

$$\sum_{j=i-1, j \geq 0}^{s-2} (-1)^{j+1} 1^{\otimes j} \otimes \phi \mu_A \otimes 1^{\otimes s-j-2}. \quad (3)$$

Main Results

In other words, the whole composition $(\phi^{[\otimes 2]} \mu^{(2)}) \circ \dots \circ (\phi^{[\otimes n-1]} \mu^{(n-1)})$ in the formula of m_n can be expressed by

$$\sum_{i_{n-1}=0}^{n-2} \left(\dots \left(\sum_{i_2=i_3-1}^1 (\phi\mu)^{(2,i_2)} \right) \dots \right) (\phi\mu)^{(n-1,i_{n-1})},$$

where $(\phi\mu)^{(k,j)} = (-1)^{j+1} 1^{\otimes j} \otimes \phi\mu_A \otimes 1^{\otimes k-j-1}$ and each addend exists whenever the corresponding index $i_k \geq 0$.

Scheme of the proof

$$m_n = (-1)^{n+1} f \mu^{(1)} \phi^{[\otimes 2]} \mu^{(2)} \dots \phi^{[\otimes n-2]} \mu^{(n-2)} \underbrace{\phi^{[\otimes n-1]} \mu^{(n-1)} g^{\otimes n}}_{k=1}$$

$$\underbrace{\hspace{10em}}_{k=2}$$

$$\underbrace{\hspace{15em}}_{k=n-2}$$

Scheme of the proof

$$\boxed{k = 1} \quad \phi^{[\otimes n-1]} \mu^{(n-1)} g^{\otimes n}$$

$$\left(\sum_{j=0}^{n-2} 1^{\otimes j} \otimes \phi \otimes (gf)^{\otimes n-j-2} \right) \circ \left(\sum_{i=0}^{n-2} \pm g^{\otimes i} \otimes \mu_A g^{\otimes 2} \otimes g^{\otimes n-i-2} \right)$$

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$$\phi g = 0$$

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$$\begin{aligned} \phi g &= 0 \\ fg &= 1 \end{aligned}$$

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$$\sum_{i=0}^{n-2} \pm 1^{\otimes i} \otimes \phi \mu_A \otimes 1^{\otimes n-i-2}$$

Scheme of the proof

$$k = 2$$

$$(\phi^{[\otimes n-2]} \mu^{(n-2)}) \circ \left(\sum_{i=0}^{n-2} \pm g^{\otimes i} \otimes \phi(\text{sth.}) \otimes g^{\otimes n-i-2} \right)$$

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Initial formulation

$$m_n = (-1)^{n+1} f \mu^{(1)} \phi^{[\otimes 2]} \mu^{(2)} \dots \phi^{[\otimes n-1]} \mu^{(n-1)} g^{\otimes n}, \quad n \geq 2.$$

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The number of basic operations can be expressed by $n + n(n-1)!^2 + \frac{(n+3)(n-2)}{4} (n-1)!^2$, so m_n is $O((n)!^2)$ in time.

First reduction

$$\phi^{[\otimes k]} \mu^{(k)} = \sum_{i=0}^{k-1} (-1)^{i+1} \mathbf{1}^{\otimes i} \otimes \phi \mu_A \otimes \mathbf{1}^{\otimes k-i-1}.$$

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$\phi^{[\otimes k]} \mu^{(k)}$: k^2 addends $\rightarrow k$ addends.
 So, now m_n is $O((n-1)!)$ in space.

First reduction

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So, now m_n is $O((n-1)!)$ in space.

Now, the number of operations is exactly

$$n + (n-1)!(2n-2),$$

so, m_n is $O((n)!)$ in time.

Second reduction

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The number of addends becomes $(n-1)! - S_n$,

$$\frac{(n-1)!}{2} < (n-1)! - S_n < (n-1)!,$$

But $(n-1)! - S_n$ is much "closer" to $\frac{(n-1)!}{2}$ than to $(n-1)!$:

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n	5	10	50	100
$((n-1)! - S_n)/(n-1)!$	0,70833	0,5637	0,51042	0,5051

Comparative table

Summing up,

	original formula		new formula	
	time	space	time	space
m_n	$O(n!^2)$	$O((n-1)!^2)$	$O(n!)$	$O((n-1)!)$